TURKISH AVIATION ACADEMY

İTÜ

Secure Multiparty Computation

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M.Sc. Program

Advanced Information Systems

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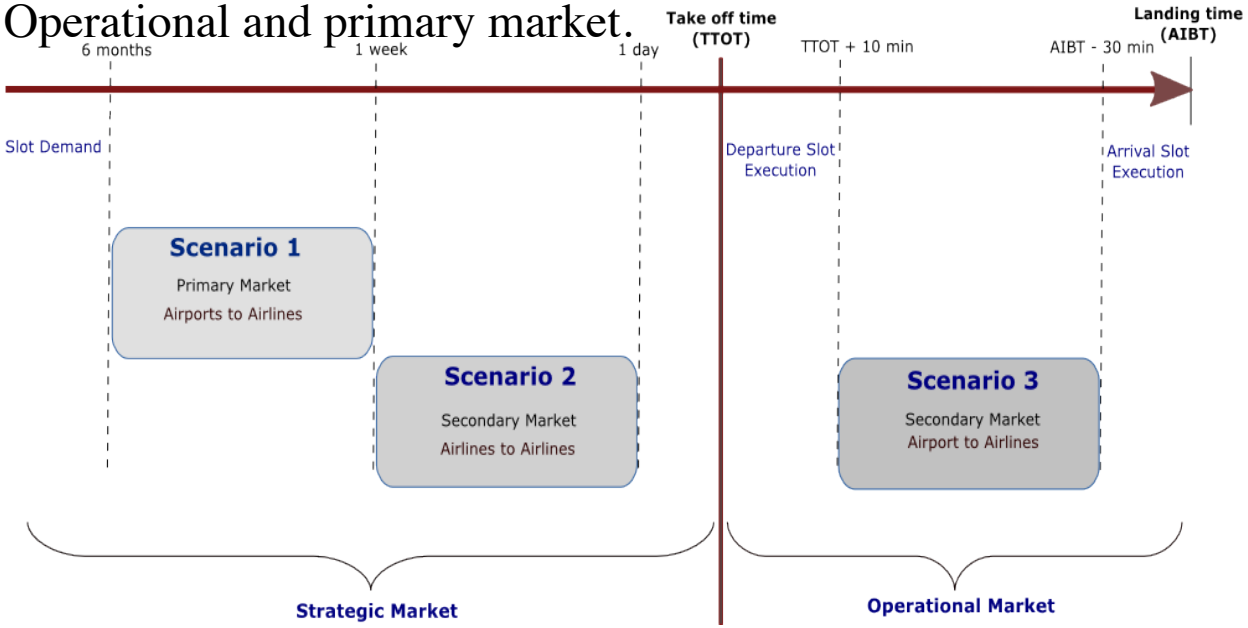
- Case scenarios:
  - **Scenario 1:** N airlines try to buy slots from an airport.
    Strategic and primary market.
  - **Scenario 2:** N airlines try to buy slots from another airline.
    Strategic and secondary market.
  - **Scenario 3:** N airlines try to buy a priority approach from an airport.
    Operational and primary market.



- General architecture

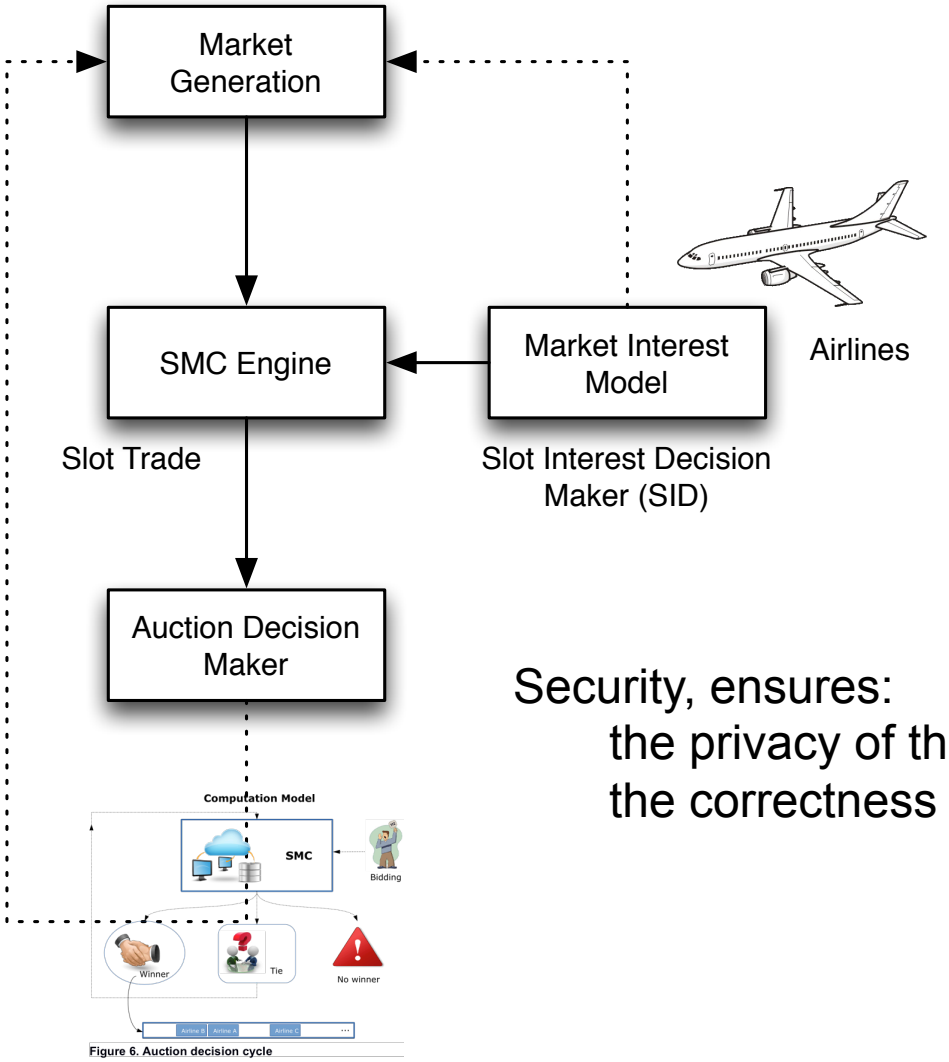


Figure 6. Auction decision cycle

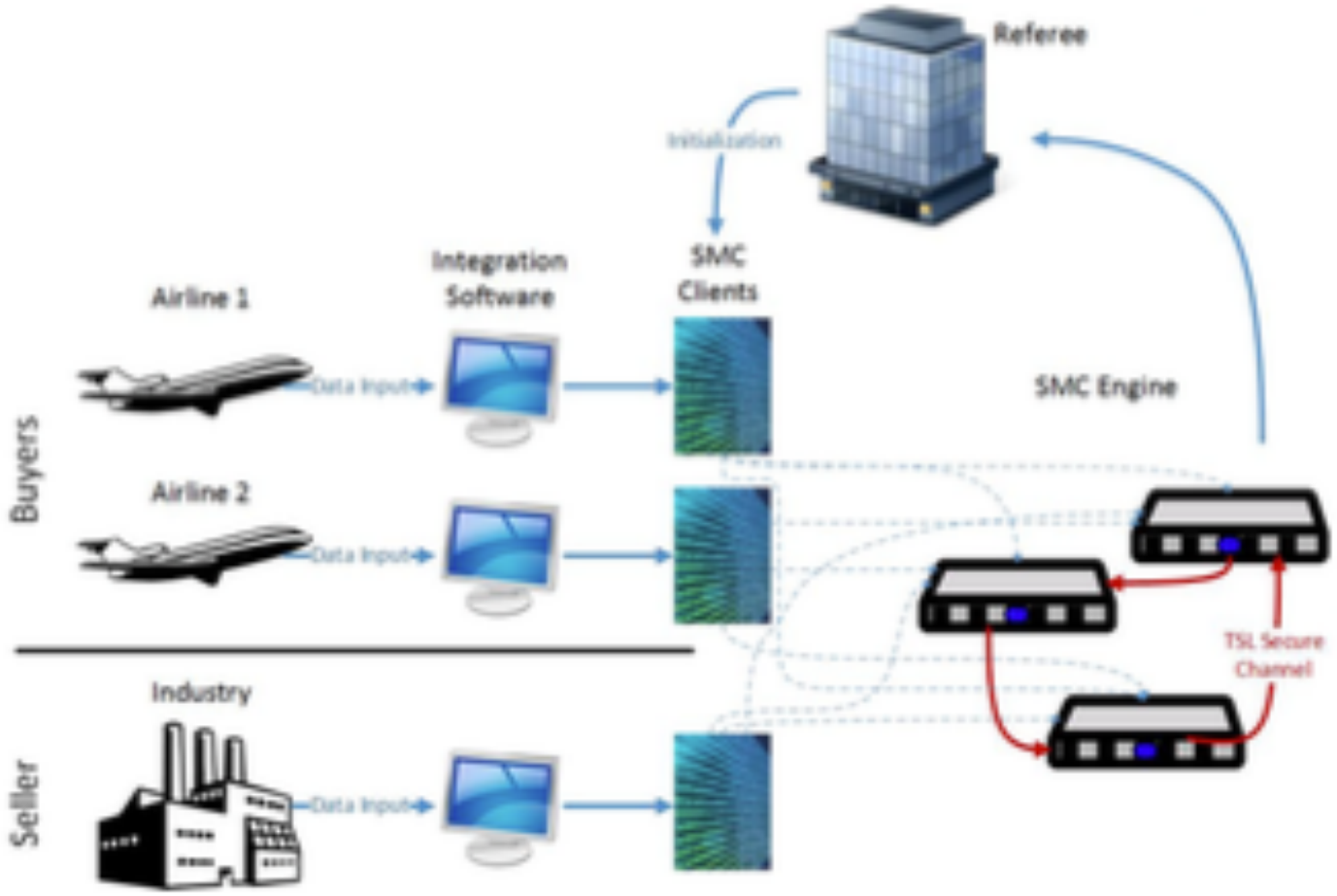
Security, ensures:

the privacy of the bidders' input and
the correctness of the computation

Slot Trading & Dynamic Landing Queues

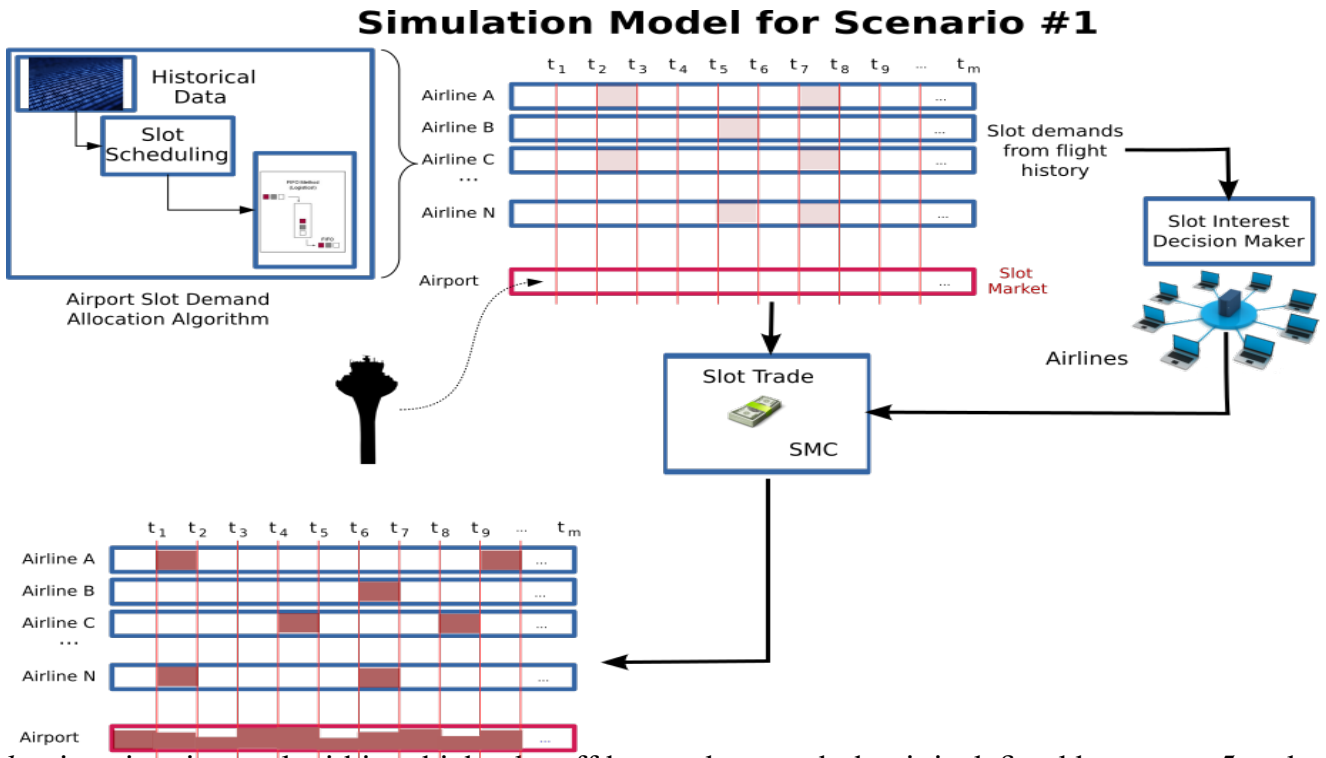
Scenario 1

Scenario 1 for Slot Trading



SCENARIO 1 FOR SLOT TRADING

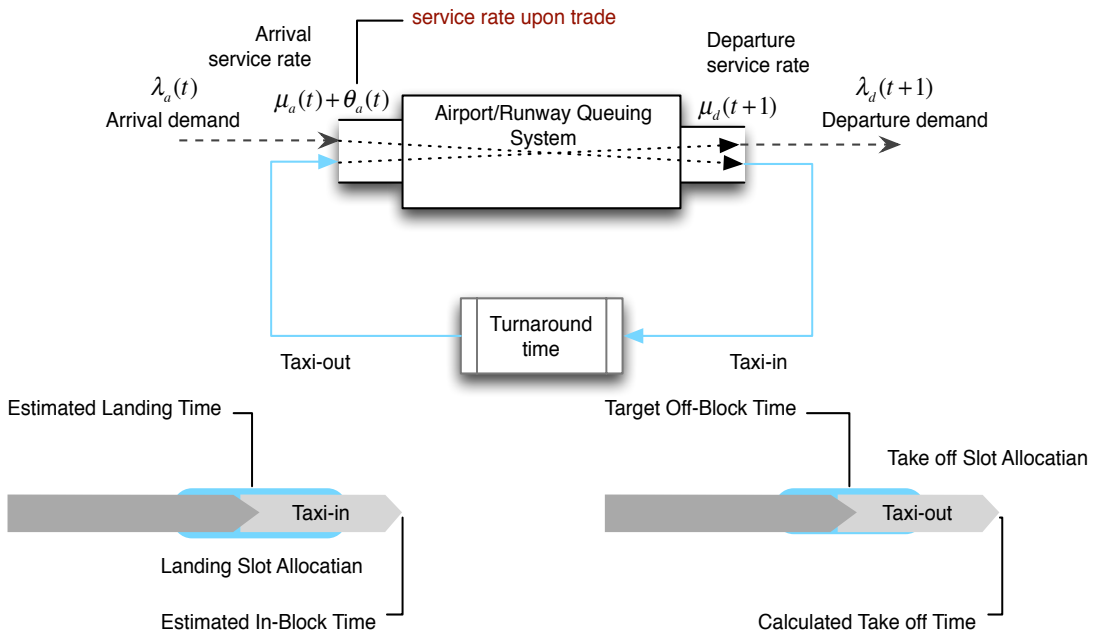
Scenario 1 focuses on initial slot trade from the airports for a long term (6 months before)



slot; is a time interval within which take off has to done such that it is defined between –5 and + 10 minutes from *Calculated Take of Time* – CTOT in Europe (EUROCONTROL 2015).

SLOT DEMAND ALLOCATION MODEL (SAM)

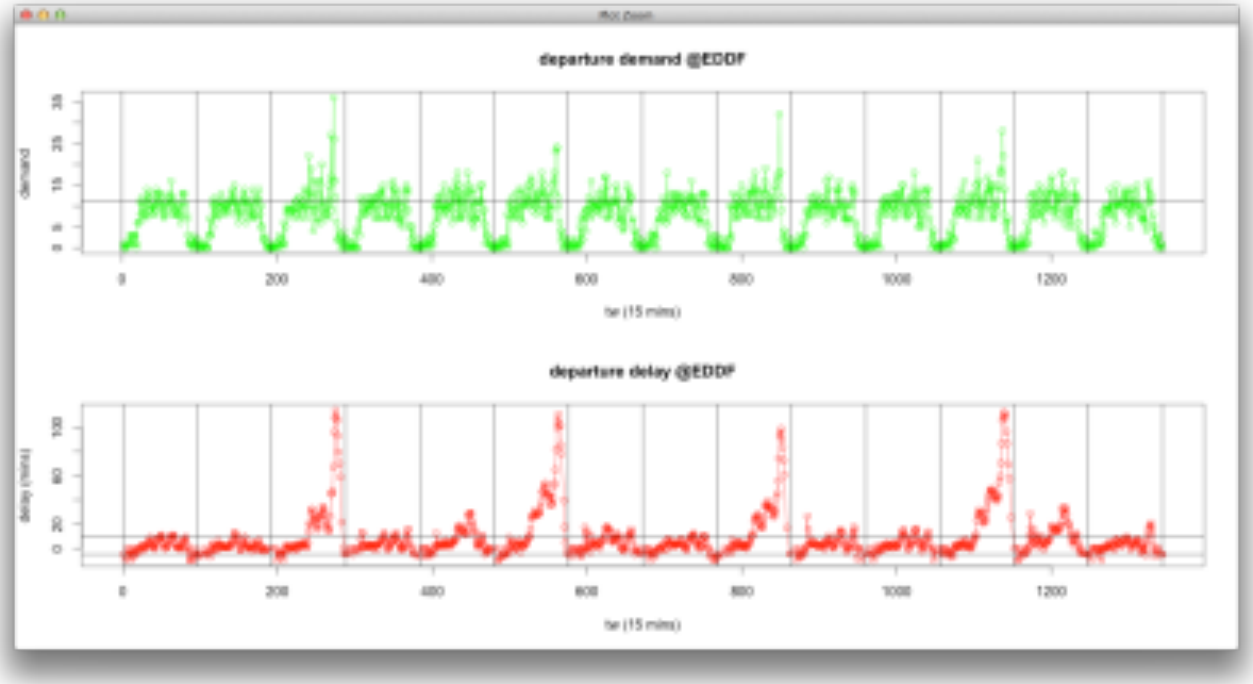
- Airport Queuing Model



– Static/Dynamics distribution of $\mu_a(t)$ to demand $\lambda_a(t)$

SLOT DEMAND ALLOCATION MODEL (SAM)

- Fill rates of all available slots vary quite a bit across airports
- Current slot capacities are based on the declared arrival and departure capacities



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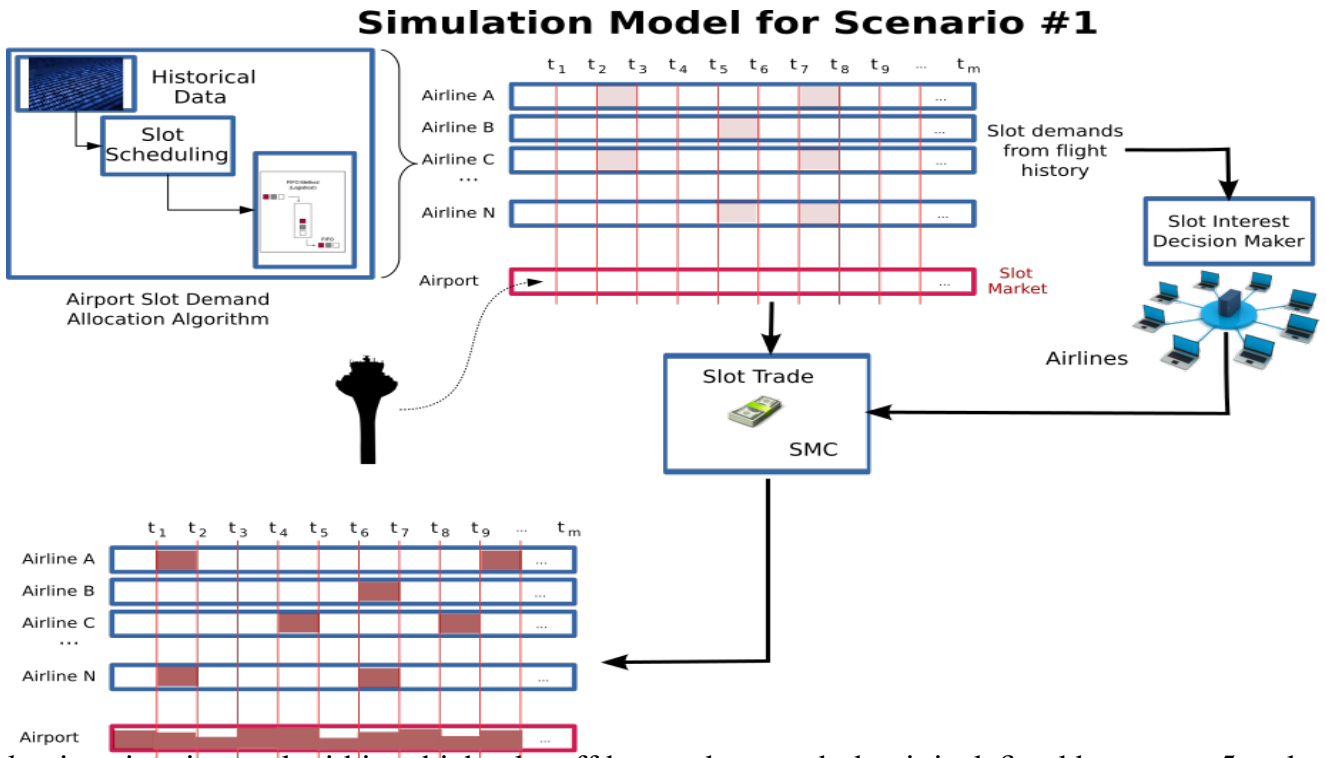

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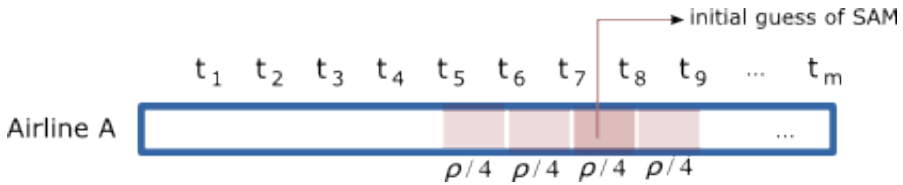
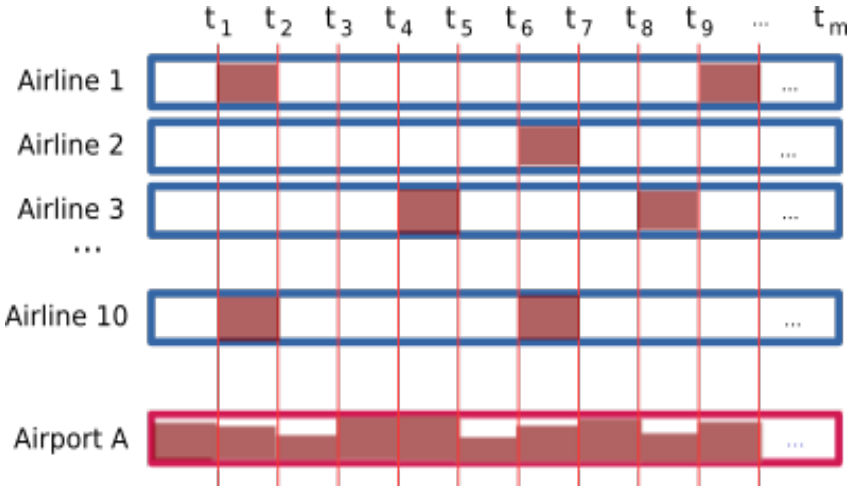
SLOT INTEREST DECISION MAKER (SID)

- *Slot Demand Allocation Model (SAM)* builds schedules for the interest of the airlines.
- Competition begins when at least two airlines want to get same slot exceeding its capacity.



SMC BID COMPARISON

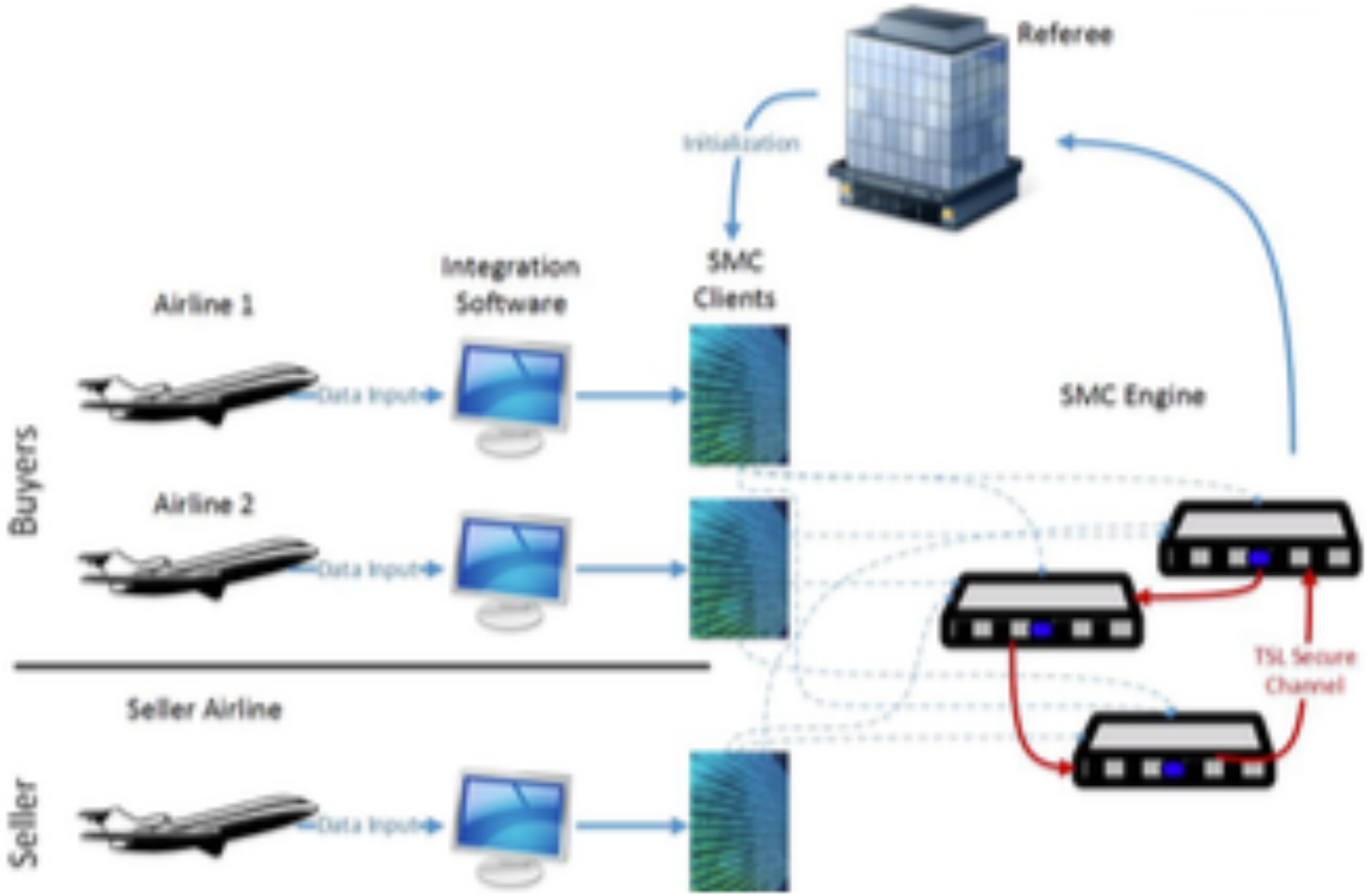
- SMC Engine collects offers.
  - If the result is a tie between two or more participants, the referee provides a notification to the participants and creates a new secure auction.
  - If the minimum price that the seller asks for has not been reached, participants are informed about this.
- Winning price and the winner is disclosed to all participants



Slot Trading & Dynamic Landing Queues

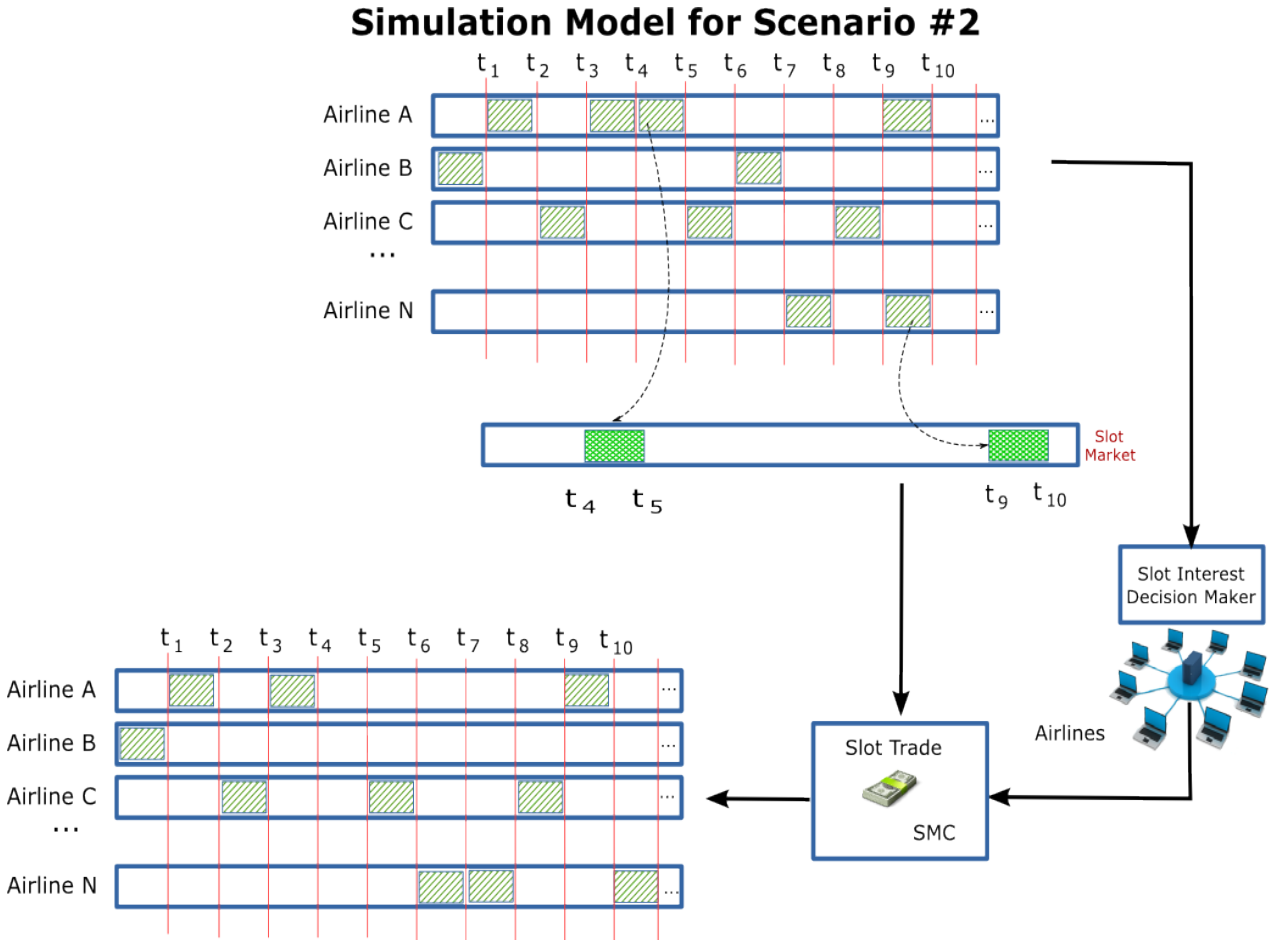
Scenario 2

Scenario 2 in Slot Trading



SCENARIO 2 IN SLOT TRADING

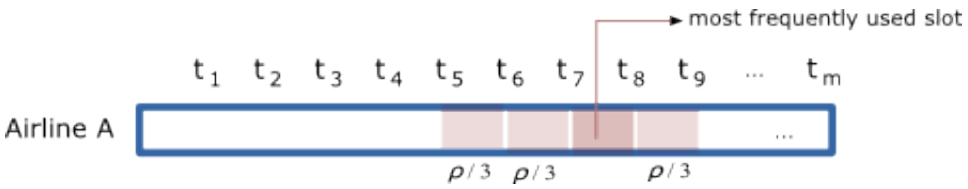
- Scenario 2 focuses on secondary strategic market in which airlines trades their slot between them.

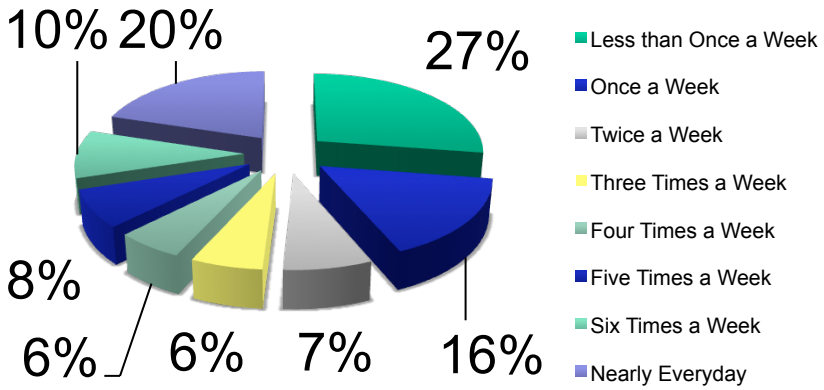


SLOT INTEREST DECISION MAKER (SID)

- *Slot Interest Decision Maker (SID)*, hypothetically generates the market and the traders by utilizing flight frequencies

  - The idea behind the concept model is that
    - airlines are interested in selling non-scheduled or least frequently used slots
    - airlines are interested in buying the slots around their most frequently used slots





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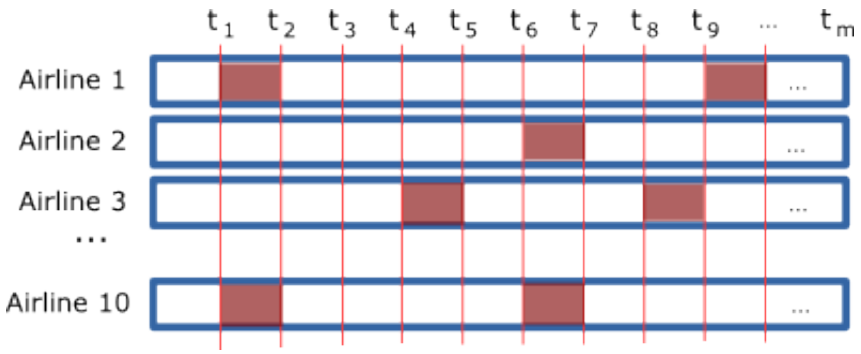
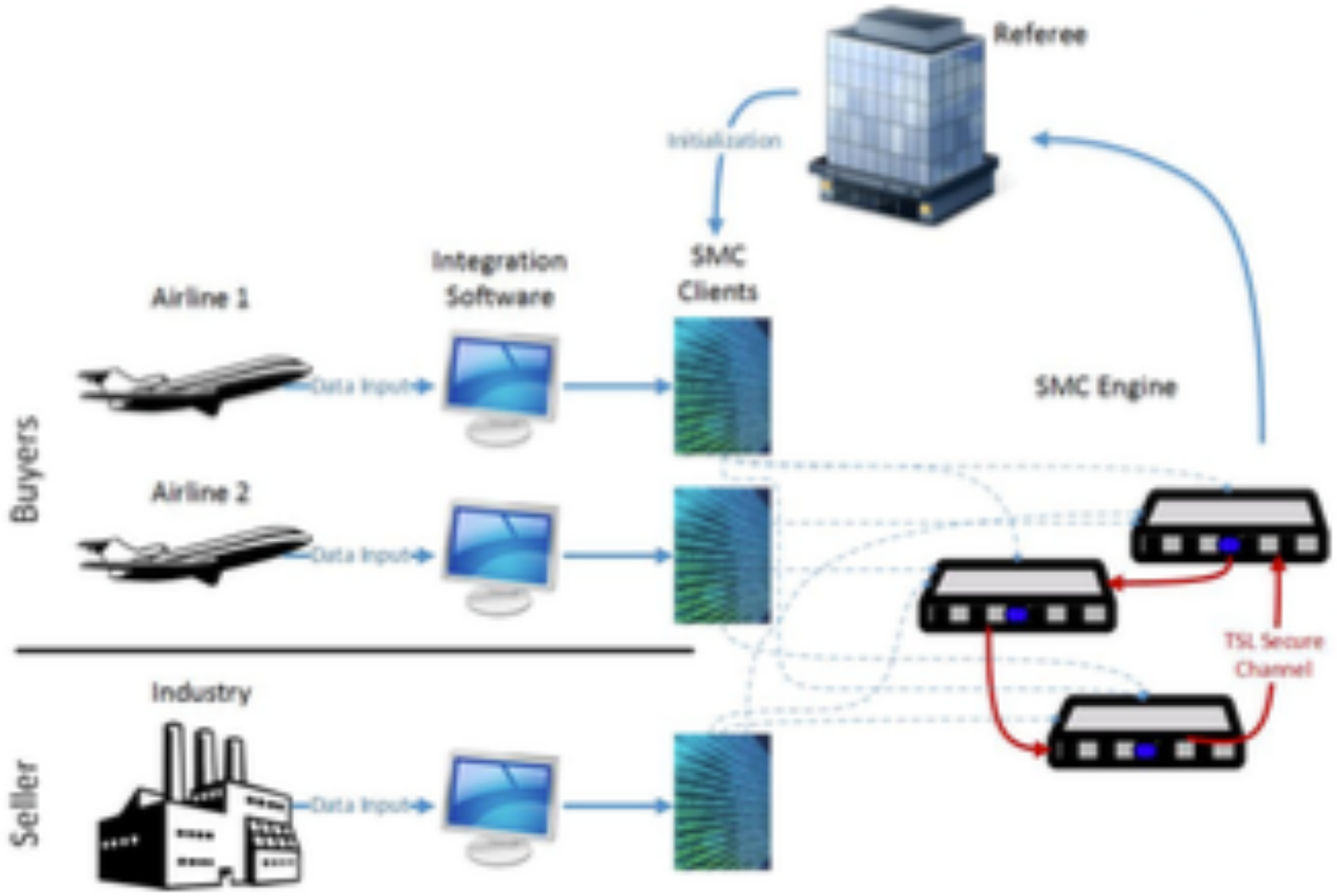


Slot Trading & Dynamic Landing Queues

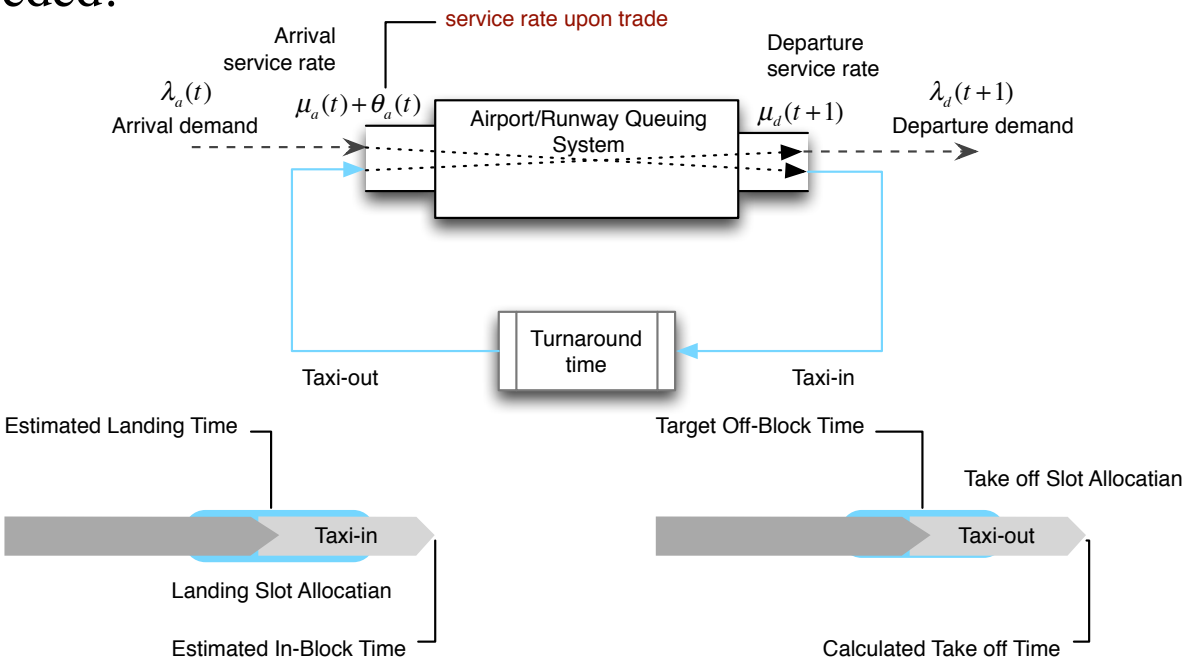
Scenario 3

Scenario 3 for Slot Trading



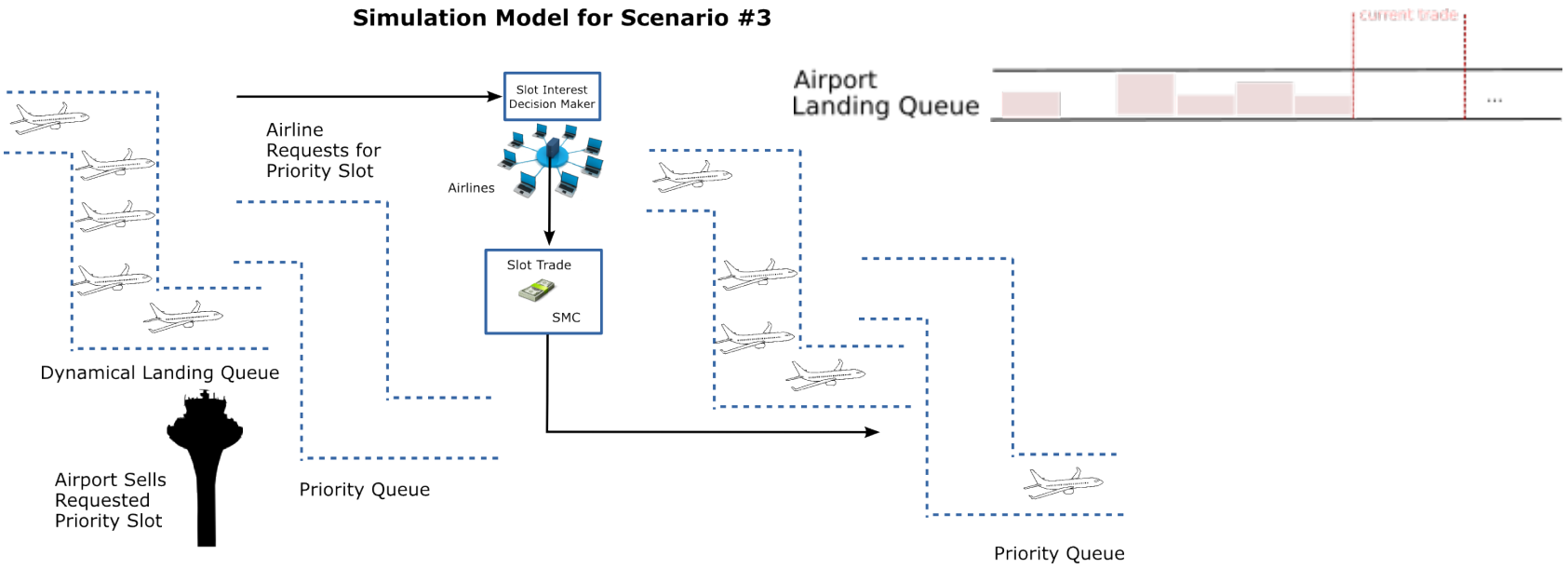
SCENARIO 3

- Scenario 3 focuses on buying a priority landing slot from an airport during approach.
  - Suppose that the airport has an additional runway with capacity $\theta_a(t)$ which is open to trade
  - This particular runway can be utilized in normal operations, however bid customers can have priority utilizing this capacity as needed.

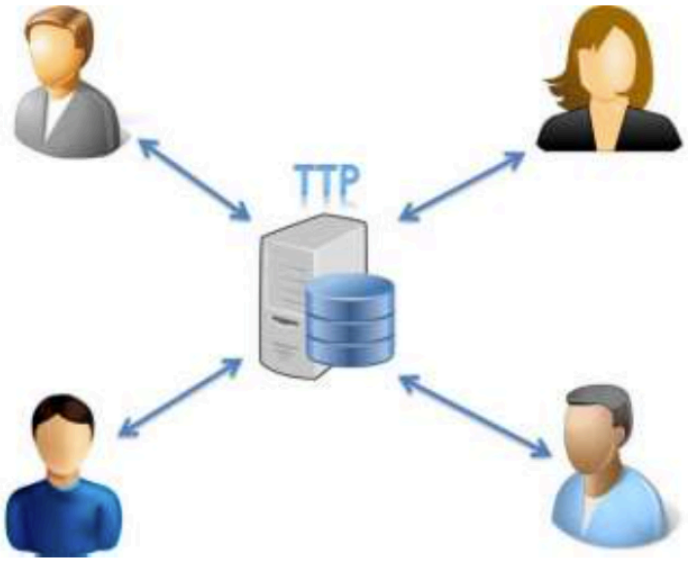


SLOT INTEREST DECISION MAKER (SID)

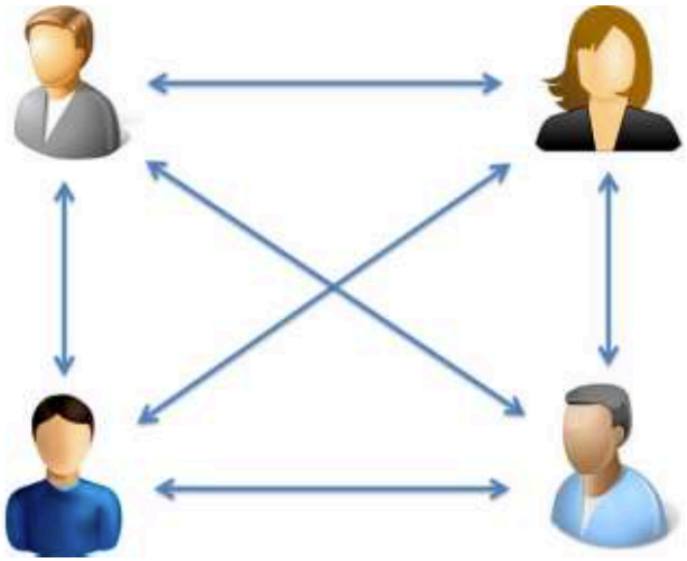
- *Slot Interest Decision Maker (SID)* utilizes current delay
  - Evaluate en-route delays for each landing aircraft.
  - If it is more than 30 minutes, airline is interested in buying a priority slot.



Secure Multiparty Computation



Trusted Third Party (TTP), Traditional Model



Secure Multiparty Computation (SMC)

SECURE SHARING?

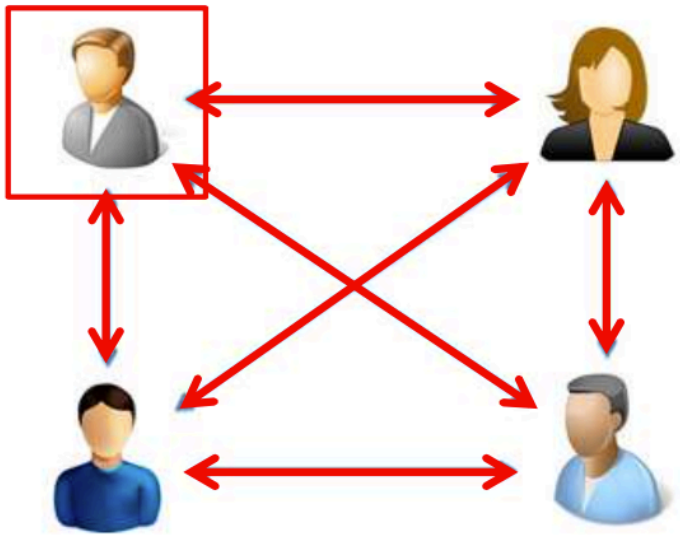
Subfield of cryptography

- Term coined by Yao (1982) in the “Millionaire problem”:

  “Two millionaires wish to know who is richer; however, none of them wants the other to find out how much his fortune is worth. How can they know who’s the richest?”

- N players wish to securely compute a given function
  - No one learns anything else than its private input and the result of the computation
- Security ensures:
  - the privacy of the player’s input
  - The correctness of the computation

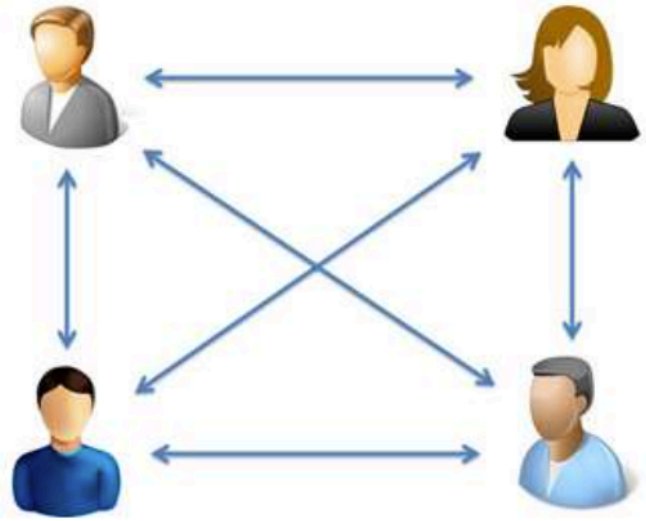
MULTIPARTY?



Party: A **participant** in the secure computation, also called a **player.**

Protocol: In general terms, it describes **how the algorithms should be used, how the players interact.**

SEMI-HONEST PARTY?



Semi-honest party (honest but curious party):

- **parties** in the computation are **corrupted** by a **passive adversary**
- **parties** will always **follow** the **protocol correctly**
- **parties** will try to learn the others private data by examining all the data they get
- the **outcome** of the **computation won't be affected**

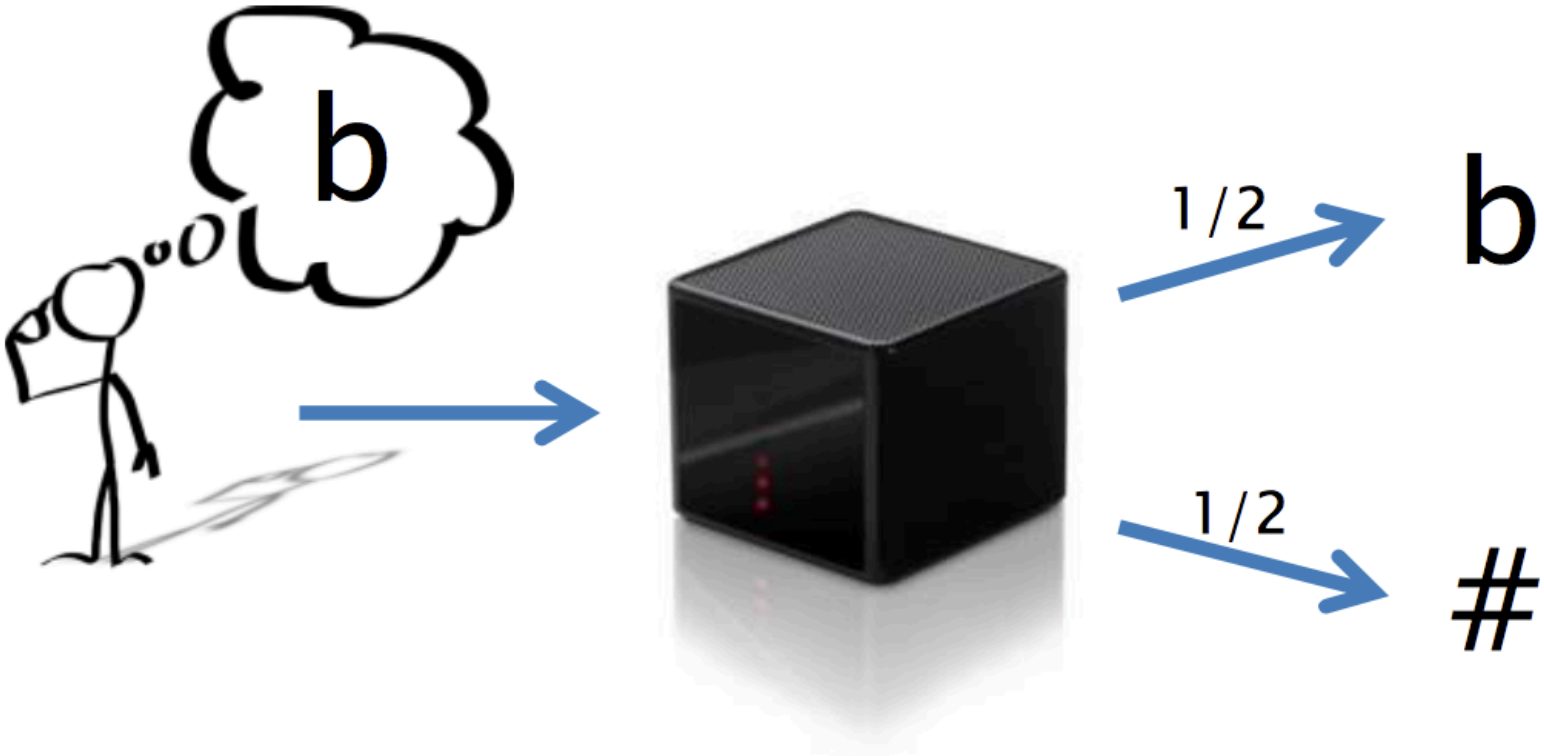
SMC Primitives

Primitive	Remarks	Usage
Oblivious Transfer (OT)	• Sends data with a ½ probability of delivery • Used as an arithmetic black box	Usually used in two-party computations
Secret Sharing (SS)	• Split the private data into shares • Unbreakable without the needed number of shares	Most widely used primitive in the construction of SMC protocols
Homomorphic Encryption (HE)	• Doesn't need to generate shares, it can perform the computation over the encrypted data directly • Very powerful, but complex to implement	Theoretical approach, too costly and complex to implement

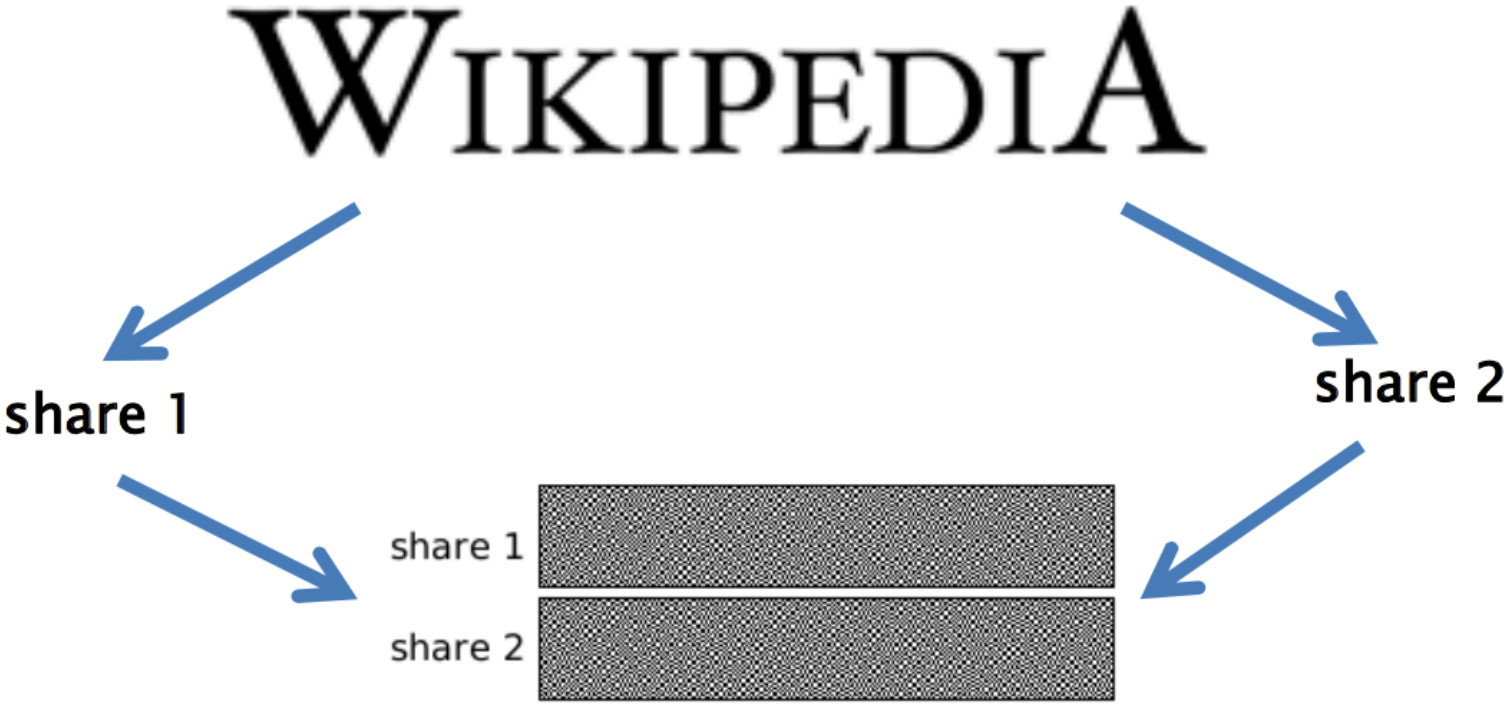
OBLIVIOUS TRANSFER

Rabin's oblivious transfer is kind of formalization of "noisy wire" communication

- Sender sends bit b into OT machine
- Machine then flips the coin, and with probability 1/2 sends b to receiver.
- Sender does not know which output receiver received



Secret Sharing



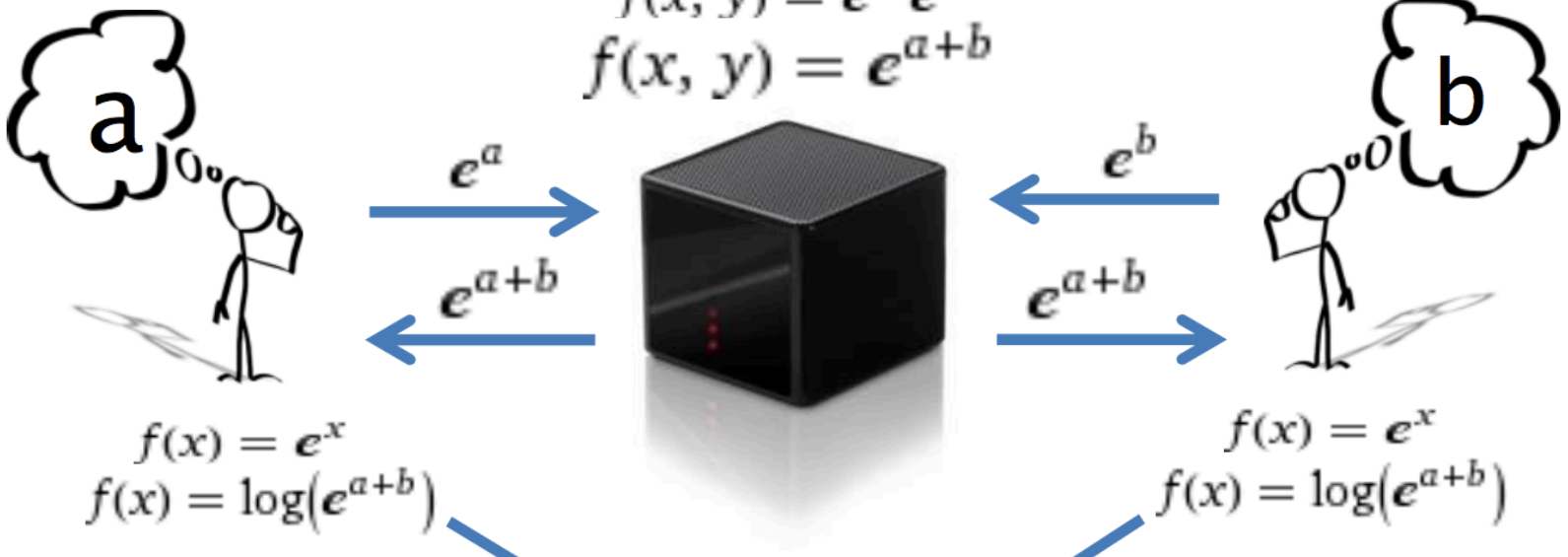
HOMOMORPHIC ENCRYPTION

HOMOMORPHIC ENCRYPTION

Let's calculate a + b!

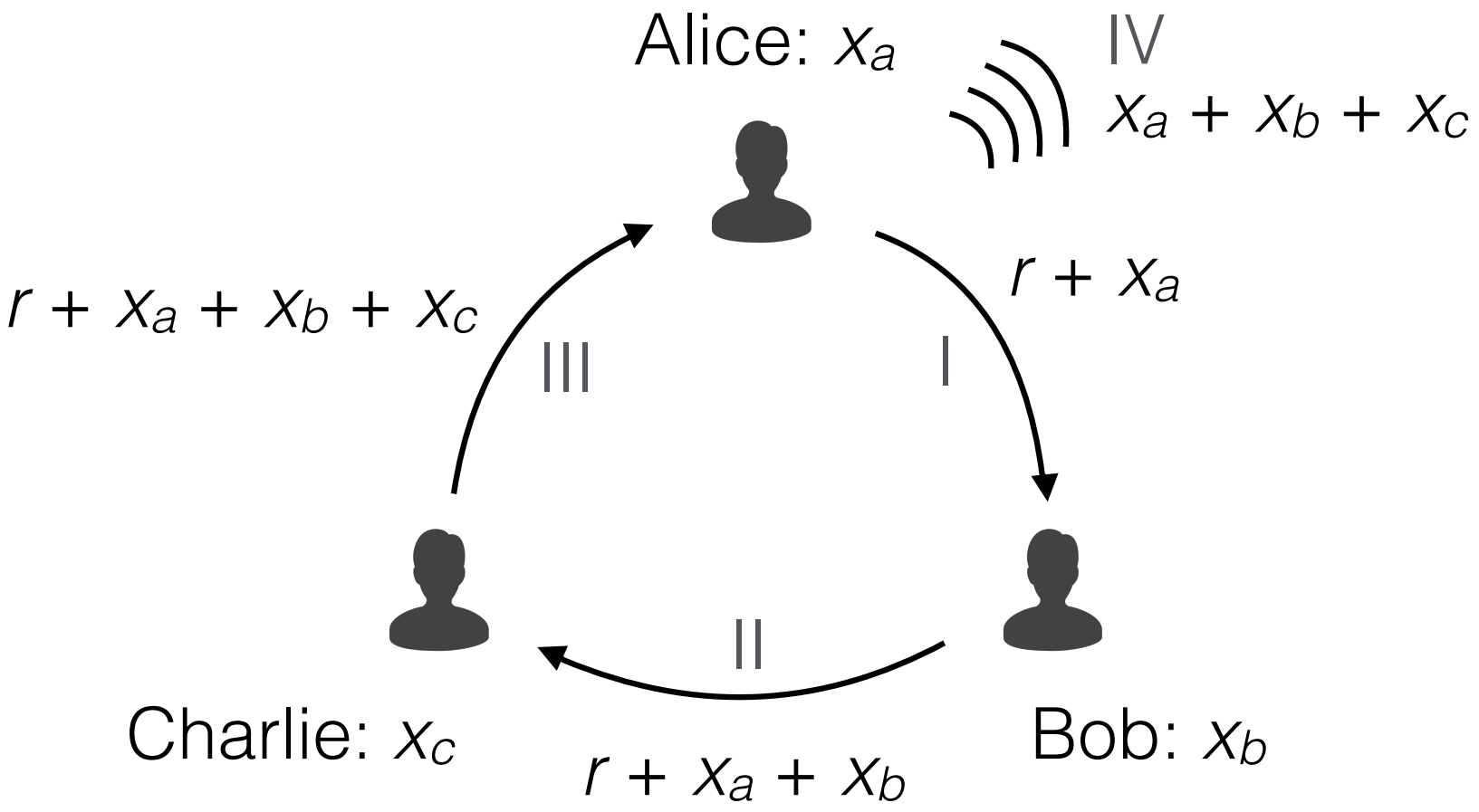
$$\exp(x+y) = (\exp x)(\exp y)$$

$$f(x, y) = x\,y$$
$$f(x, y) = e^{a}\,e^{b}$$
$$f(x, y) = e^{a+b}$$



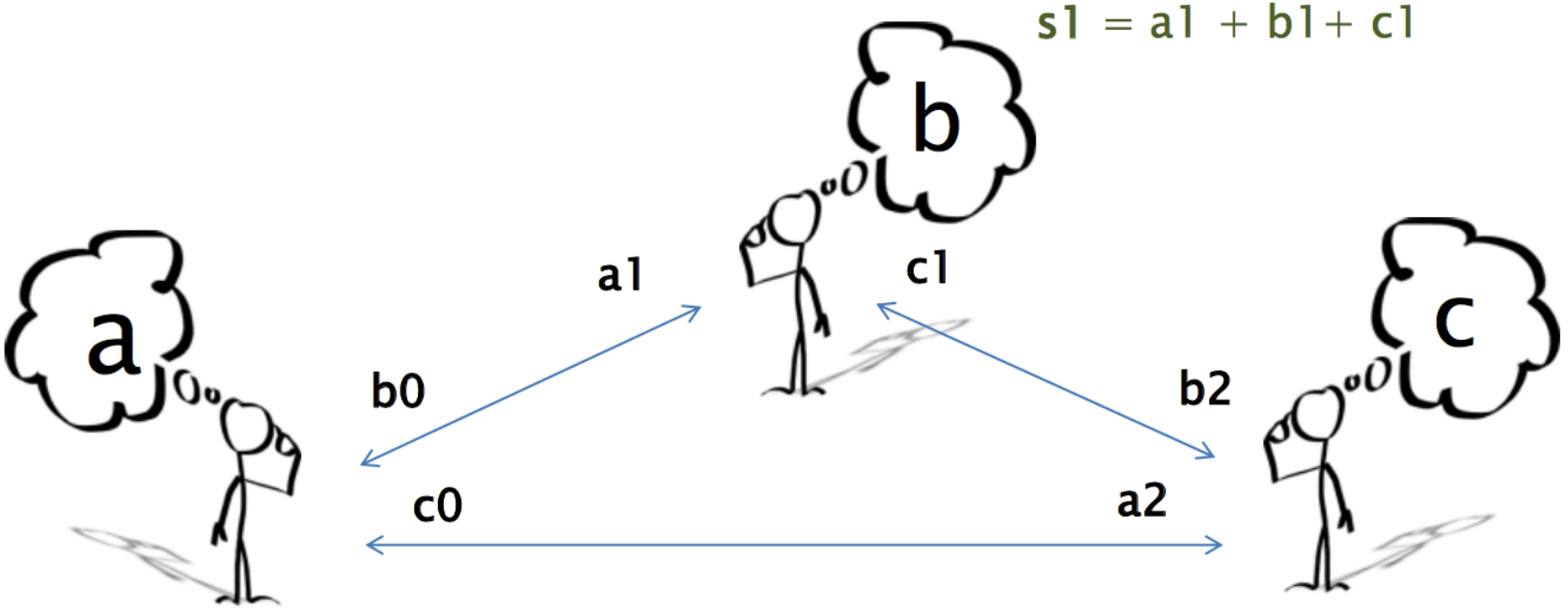
$$a + b = f(x)$$

Secret Sharing (I)



SECRET SHARING (II)

Let's *securely* calculate a + b + c!



$$s0 + s1 + s2 = a + b + c$$

SECURE BIDDING

units	price/unit
1	20
2	9
3	5,5

Market clearance price

units	price/unit
1	18
2	9
3	5

THIRD PARTY SMC ENGINE

Company1

Company 2

Sector benchmarking:

- Statistical Data
- Strategic Ranking
- Anonymity